



Detecting Covert Communications on the Internet: Some Challenges and Solutions

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Covert Channel Definition

Any communication channel that can be exploited by a process to transfer information in a manner that violates the system security policy.

--- U.S.D.O.D. 1985, Trusted computer system evaluation criteria.

Example Internet Covert Channels

- TCP/IP Protocol
 - » Unused header fields.
 - » Encoding information in sequence numbers.
- Timing Channel
 - » Encode covert info. in the rate at which jobs are sent to a time-shared server.
 - » Measuring response times to jobs gives noisy version of message.
- Digital media (image, video...) on the web

Example Covert Messages

- Spy programs
- Company proprietary information.
- Computer virus...

Issues in Intercepting Covert Messages on Internet (I)

- What to look for in the Internet?
- Where to look? How to identify web links that could potentially contain covert messages?
 - » Side information may be available.
 - » No side information at all.
 - random search may be futile.
 - metrics used by current web search engines may not work: e.g., popular sites.

Issues in Intercepting Covert Messages on Internet (II)

- Message carrying website may have not a public link
 - » Use http traffic request in the back bone to identify these hidden links?
- Websites are created, moved, and destroyed randomly on a daily basis
 - » Continually monitor websites of interest?
 - » How often to monitor?
- A web-page like e-bay could contain thousands of images
 - » Efficient search techniques.

Some Approaches

- Candidate websites for investigation could be chosen as follows:
 - » External info. such as email trace, phone tapping, etc.
 - » Eliminate certain sites such as .mil, .gov...
 - » Past history.
 - » Religious cult's website?
 - » Websites of groups with radically politically opposed views?
 - » Info. from network forensic tools.

Steganography Covert Channel Requirement

- Maximize stealth
 - » Detection via steganalysis is "difficult."
 - » Perceptually transparent
- Maximize capacity
 - » Maximum embedded message length such that steganalysis detection is "difficult."
- Efficient encoding/decoding

Intercepting Steganographic Messages (I)

- **Steganalysis**
 - » Analyze digital data to determine presence of secret messages.
- **Passive Steganalysis**
 - » Steganalyst/hacker/interceptor tries to find if a secret message is present.
 - » Identify the embedding algorithm/software used.
 - » Removal of secret message is not an aim.
 - » Little or no *a priori* information available.
- **Active Steganalysis**
 - » Estimate the secret key, message length, etc.
 - » Estimate the secret message (grand goal!).

Intercepting Steganographic Messages (II)

- Theoretical issues.

- » What must a steganalysis algorithm look for?
- » What are the "give aways" in current published steganography algorithms/software?
- » What is the minimum message length that can be detected?
- » What about false alarm and miss probabilities?
- » Mathematical tools from probability and statistics.

- Scalability.

- » Investigating every web site is not possible.
- » How often web sites are to be investigated?
- » Can we identify "high risk" sites in some sense?
- » Number of possible embedding algorithms could be large.
- » Message sizes could be small.

Intercepting Steganographic Messages (III)

- Is steganalysis realistic?
 - » Current approaches seem to be extreme:
 - tuned to work for one particular embedding algorithm or use large training data set.
 - » What if the embedding algorithm is not published publicly?
 - » Need: steganalysis that works for a "class" of embedding algorithms.

Steganalysis Current Trends (I)

- Classifier/statistical learning based
 - » Train steganalysis classifiers on large training sets
 - » Use host data features.
- Pros
 - » Well understood classifier theory
 - » Works reasonably well in practice
- Cons
 - » May not work well for data that are significantly different from training set
 - » Overfitting problems
 - » How to choose training set? How large a training set? Mostly heuristics involved here.

Steganalysis Current Trends (II)

- Blind statistical system identification based
 - » Use individual host data features.
 - » No training set.
- Pros
 - » Sound theoretical analysis possible.
 - » Covert message extraction demonstrated.
- Cons
 - » Stochastic non-stationarity of digital data, e.g., images.
 - » New tricks needed to make it work in practice.

Optimal Web Search for Covert Message: A Mathematical Model

- Let,
 - » Total number of web sites to be searched = W .
 - » $P(j) = \text{Pr}(\text{website "j" contains the covert message})$.
 - » $b(j,t) = \text{Pr}(\text{detecting covert message after spending } t \text{ time units in site } j \mid \text{message in site } j)$.
 - » $C(j,f(t)) = \text{cost of searching site } j \text{ with a time/resource allocation of } f(t)$.
 - » $P(f) = \sum_j P(j)b(j,f(t))$ is average probability of finding covert message.
 - » $C(f) = \sum_j C(j,f(t))$ is total cost.

Possible Scenarios

- Case 1: $\{P(j); j=1,2,\dots,W\}$ completely known.
 - » Subjectively chosen.
- Case 2: Only ordering of probabilities known, i.e., $P(1) > P(2) > \dots > P(W)$
 - » More realistic.
- Case 3: $\{P(j)\}$ completely unknown.
 - » When no side info. available.

- Web search strategy for Case 1:
 - » $P(f^*) = \max P(f)$ subject to $C(f) < T$.
- $f^*(t)$ is the optimal allocation of time to search each of the web sites.
- Suppose $b(j,t) = 1 - e^{-t}$, $t > 0$ then, $f^*(t)$:

$$t_j = \max(0, \ln(p_j/K)); j=1,2,\dots,W$$

where $K = [\prod p_j]^{1/W} e^{-T/W}$

$$1 - e^{-T/M} \leq P(f^*) \leq 1 - T [\prod p_j]^{1/W} e^{-T/W}$$

- » For a desired covert message detection accuracy, bound on total required resource T can be computed.

Web Search Strategies (I)

- Let probability of detecting covert website/channel = q
- Probability of success statistically independent from one attempt to another.
- Possible search strategies:
 - » Co-ordinated strategy.
 - » Randomized strategy.

Web Search Strategies (II)

- Co-ordinated search
 - » Results of previous search results stored. That is, "memory" is built into searching.
 - » Web links previously searched, images previously investigated, etc. are stored.
 - » Avoid these links/data in future searches.
 - » Pros: Optimal strategy because of the memory.
 - » Cons: Large storage needed, cached data could become outdated...

Web Search Strategies (III)

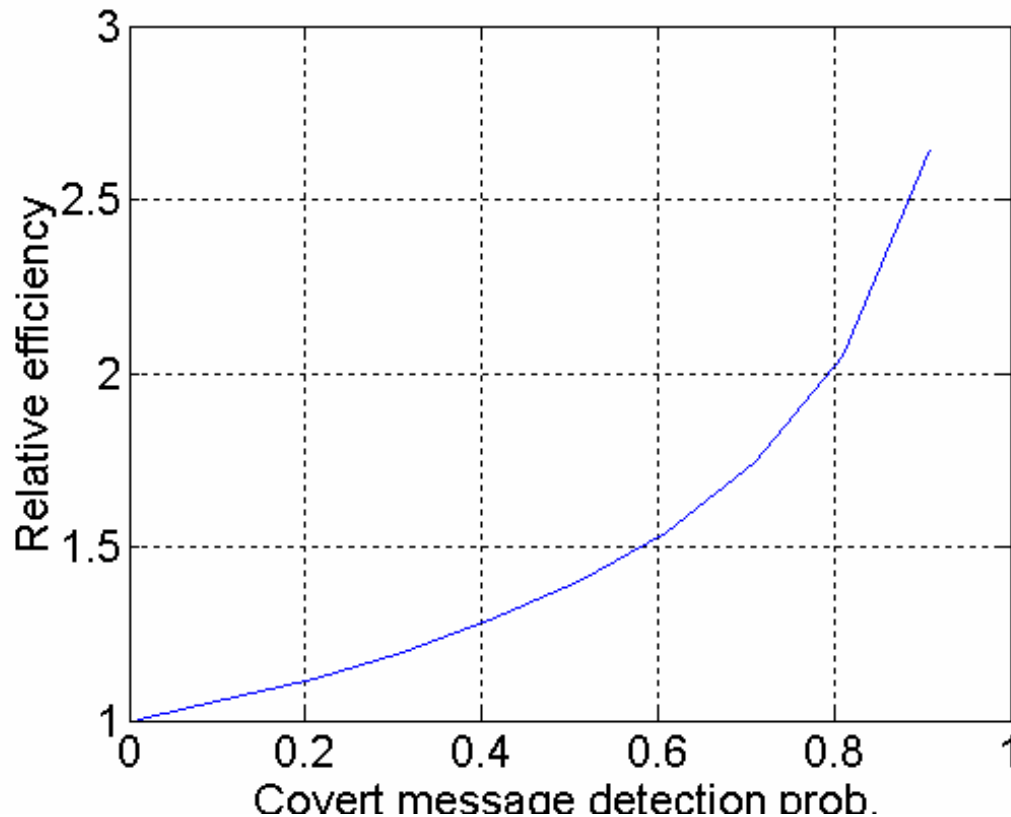
- Randomized search
 - » No memory built.
 - » Search web sites randomly.
 - » Pros: Large storage is not needed.
 - » Cons: Does not exploit memory.

Search Relative Efficiency (I)

- N_r = no. of times a web site is searched for detecting a covert message using randomized search.
- N_c = no. of times a web site is searched for detecting a covert message using co-ordinated search.
- d_r = prob. Of detecting covert message using randomize search.
- d_c = prob. Of detecting covert message using randomize search.

Search Relative Efficiency (II)

- If $d_r = d_c$ then, rel. eff. of the random search w.r.t. co-ord. = $-\ln(1-q)/q$.



Key Observation

- If covert message detection reliability is low, then co-ordinated and randomized searches are nearly equally efficient.

Additional Information

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