# Pool Tag Quick Scanning for Windows Memory Analysis

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# Who are any of us really?

- Senior Research Developer @BlackBag Tech
  - find cool artifacts, figure out how to parse them
  - develop new techniques
  - get them into our tools
- PhD from UNO in CS (2009)
  - research focus on efficient digital forensics
- Also done: DF practice, PenTesting, Malware Analysis
- FOSS dev: Scalpel, Registry Decoder, Spotlight Inspector, DAMM (built on top of Volatility)
- Organizer BSidesNOLA April 16 in New Orleans
  - come on over and I'll buy the beers





## The Problem

- Memory forensics coming into wider use
- Applications for DF
  - crypto, cached data, volatile system state
- And IR
  - malware, intrusion detection
- Just like disk, memory sizes are increasing rapidly
  - newer Windows systems max out at ~4TB
- Some memory analysis relies on scanning
  - like file carving but for in-memory structures
- I (and likely you) want everything to be faster
  - without loss of ... anything
  - especially in IR Land





# The (Basic) Solution

- Generally, the most important things we scan for are kernel structures
  - e.g., \_EPROCESS process descriptors
- These things exist in kernel memory
- Kernel memory divided into a set of pools
- Many of the things we care about are only allocated from specific pools
  - a much smaller scanning space





# Memory Pools

- Dynamically sized (heaps)
- Kernel allocations in system address range
  - kernel address space
  - mapped into every process
- Paged pool: can be paged out to disk
- Non-paged pool: cannot be paged out to disk
  - so guaranteed to be in a memory image
  - kernel structures (processes, network stuff)
  - drivers
  - observed as small as 64MiB (allocated)
- Begin with a \_POOL\_HEADER structure





```
nt!_POOL_HEADER
```

struct \_POOL\_HEADER, 9 elements, 0x10 bytes

+0x000 PreviousSize : Bitfield Pos 0, 8 Bits

+0x000 PoolIndex : Bitfield Pos 8, 8 Bits

+0x000 BlockSize : Bitfield Pos 16, 8 Bits

+0x000 PoolType : Bitfield Pos 24, 8 Bits

+0x000 Ulong1 : Uint4B

+0x004 PoolTag : Uint4B

+0x008 ProcessBilled : Ptr64 to struct \_EPROCESS

+0x008 AllocatorBackTraceIndex : Uint2B

+0x00a PoolTagHash : Uint2B

**BlockSize**: size of allocation\*

PoolType: paged pool, non-paged pool

PoolTag: 4 byte marker for this allocation type





```
ntdll!_POOL_TYPE
Enum _POOL_TYPE, 15 total enums
   NonPagedPool = OnO
   PagedPool = On1
   NonPagedPoolMustSucceed = On2
   DontUseThisType = 0n3
   NonPagedPoolCacheAligned = 0n4
   PagedPoolCacheAligned = 0n5
   NonPagedPoolCacheAlignedMustS = 0n6
   MaxPoolType = 0n7
   NonPagedPoolSession = 0n32
   PagedPoolSession = 0n33
   NonPagedPoolMustSucceedSession = 0n34
   DontUseThisTypeSession = 0n35
   NonPagedPoolCacheAlignedSession = On36
   PagedPoolCacheAlignedSession = 0n37
   NonPagedPoolCacheAlignedMustSSession = 0n38
```





Purpose	Pool Tag	
Driver Object	Driv	
File Object	File	
Kernel Module	$\operatorname{MmLd}$	
Logon Session	SeLs	
Process	Proc	
Registry Hive	CM10	
TCP Endpoint	$_{\text{TcpE}}$	
TCP Listener	$\mathrm{TcpL}$	
Thread	Thre	
UDP Endpoint	UdpA	





# Big (Large) Page Pool

Allocations over a certain size (~page size\*) are made from the Big Page Pool
Info about allocations at nt!PoolBigPageTable

```
struct _POOL_TRACKER_BIG_PAGES, 4 elements, 0x18 bytes
+0x000 Va : Ptr64 to Void
+0x008 Key : Uint4B
+0x00c PoolType : Uint4B
+0x010 NumberOfBytes : Uint8B
```

Va: virtual address of the allocation

Key: pool tag

PoolType: type

NumberOfBytes: size of allocation





# Pool Tag Scanning

- Pool tags are handy for scanning through entire memory image
  - analogous to a file header
  - at least for smaller allocation sizes
- Also like file headers, need further verification to reduce false positives
  - use known constraints for structure type
  - other nearby structures like \_OBJECT\_HEADER
- What about Big Page Allocations?
  - pool tag stored at nt!PoolBigPageTable, not with the allocation itself
  - just enumerate the table





# Pool Tag Quick Scanning

- Crux: We know allocations for key kernel structures come from specific pools
  - non-paged pool
  - big page pool
- For non-paged pool, kernel keeps a VA allocation bitmap
  - what VAs are mapped to physical pages
- PTQS Process
  - get virtual address range of non-paged pool and use VA allocation bitmap to find those mapped physical pages
  - use big page table to find allocations backed by physical pages
  - use VAs/page tables to build range of physical pages to scan
  - scan only these pages
- Does it work?
  - Glad you asked.





## Base Test Setup

- We are currently developing a new memory analysis framework (topic of coming paper)
- Developed two plugins to search for \_EPROCESS allocations
  - psscan to exhaustively search physical memory
  - psquickscan to use the PTQS technique
- Ran a series of tests for accuracy, speed, etc.
  - Hardware: mid-2014 2.8 GHz MacBook Pro with 16 GiB RAM
  - Note: all times are average of 10 runs with highest and lowest removed





#### Scenario 1: Accuracy

- Win7SP1x64 16 GiB memory image
- Compare our psscan and psquickscan
- Compare to Volatility and Rekall

Plugin	Туре	Avg. Time	Running	Terminated	Prior Boot	Duplicate <sup>4</sup>
psquickscan	Virtual	0.129s	128	21	0	0
psscan	Physical	15.584s	128	22	15	43
psscan (Rekall)	Physical	35.967s	128	22	15	43
psscan (Volatility)	Physical	25.448s	128	21	15	43

#### Notes

- All scan types found the same number of running processes
- Two anomalies in the number of terminated processes found
- psquickscan reported reading only 80 MiB of the image





### Scenario 2: Speed

- Memory images across multiple OSs, and RAM sizes
- Compare our psscan and psquickscan

OS Version	Plugin	Data Scanned	RAM Size	Avg. Time	Running	Terminated	Duplicate
Vista SP0	psquickscan	38 MiB	1 GiB	0.083s	46	2	15
Vista SP0	psscan	$1\mathrm{GiB}$	1 GiB	0.356s	46	2	15
Vista SP1	psquickscan	$60\mathrm{MiB}$	$1\mathrm{GiB}$	0.073s	48	0	0
Vista SP1	psscan	$1\mathrm{GiB}$	1 GiB	$0.400 \mathrm{s}$	48	0	0
Vista SP2	psquickscan	76 MiB	1 GiB	0.236s	50	1	0
Vista SP2	psscan	$1\mathrm{GiB}$	$1\mathrm{GiB}$	$0.547 \mathrm{s}$	50	1	11
7 SP0	psquickscan	64 MiB	$2\mathrm{GiB}$	0.075s	43	4	0
7 SP0	psscan	$2\mathrm{GiB}$	$2\mathrm{GiB}$	0.712s	43	6	4
7 SP1	psquickscan	64 MiB	$2\mathrm{GiB}$	0.075s	50	5	0
7 SP1	psscan	$2\mathrm{GiB}$	$2\mathrm{GiB}$	0.691s	50	5	0
8	psquickscan	$44\mathrm{MiB}$	$4\mathrm{GiB}$	0.054s	36	3	0
8	psscan	$4\mathrm{GiB}$	$4\mathrm{GiB}$	1.433s	36	3	0
8.1	psquickscan	$244\mathrm{MiB}$	8 GiB	$0.170 \mathrm{s}$	45	0	0
8.1	psscan	$8\mathrm{GiB}$	$8\mathrm{GiB}$	2.977s	45	0	0

#### Notes

About an order of magnitude speedup

typo in paper!





#### Scenario 3: Network Data Transfer

- Use F-Response to mount RAM over network (gigabit)
- Compare our psscan and psquickscan

RAM Size	Plugin	Scanned	Time	Transferred
$2\mathrm{GiB}$	psquickscan	$102\mathrm{MiB}$	9.489s	116.115 MiB
$2\mathrm{GiB}$	psscan	$2\mathrm{GiB}$	28.132s	$2.014\mathrm{GiB}$
4 GiB	psquickscan	$122\mathrm{MiB}$	9.640s	177.367 MiB
4 GiB	psscan	$4\mathrm{GiB}$	56.971s	$4.027\mathrm{GiB}$
8 GiB	psquickscan	$246\mathrm{MiB}$	15.360s	299.648 MiB
8 GiB	psscan	8 GiB	$3\mathrm{m}26.449\mathrm{s}$	$8.132\mathrm{GiB}$

#### Notes

- Data transferred just greater than data scanned
- Slower networks will just make the wait more frustrating





### Scenario 4: Large Memory Image

- Test with significantly larger memory image
- Compare our psscan and psquickscan
- Compare to Volatility and Rekall

Plugin	Data Scanned	Avg. Time
psquickscan	$5.76\mathrm{GiB}$	5.797s
psscan	$192\mathrm{GiB}$	$3 \mathrm{m} 8.421 \mathrm{s}$
psscan (Rekall)	$192\mathrm{GiB}$	$6\mathrm{m}7.207\mathrm{s}$
psscan (Volatility)	$192\mathrm{GiB}$	4m42.412s

#### **Notes**

- About 2 orders of magnitude speedup versus other methods
- psscan linear in RAM size, not psquickscan





## A Note on Limitations

- Our limitations are inherent to scanning in virtual address space
- Starting in Windows 10 Microsoft obfuscates
   \_OBJECT\_HEADERs using the VA of the
   allocation
- Must scan in kernel's virtual address space
- tl;dr Existing tools may have the same limitations as us starting with Windows 10





## Conclusions

- New technique: limit pool tag scanning to pools where allocations for these objects are made
- Significantly more efficient
  - time: order of magnitude+ speedup
  - network bandwidth
- Minimal loss of accuracy
  - no processes from previous boot
  - terminated processes in deallocated pages not found
  - we'd have these limitations in Windows 10+ anyway





## Future Work

- More testing of pool sizes with different workloads
- Quantify the incidence of objects in deallocated pages
- Find a way to scan a subset of deallocated pages that might hold fun stuff





## Questions?

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